

# Enhancing Energy Efficiency in Sensor/Ad-Hoc Networks Through Dynamic Sleep Scheduling

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## ENHANCING ENERGY EFFICIENCY IN SENSOR/AD-HOC NETWORKS THROUGH DYNAMIC SLEEP SCHEDULING

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#### Abstract

Enhancing energy efficiency is a pivotal concern in sensor/ad-hoc networks, where devices are often constrained by limited power sources. Dynamic sleep scheduling emerges as a promising strategy to mitigate energy wastage and prolong network longevity. This approach involves orchestrating nodes to periodically transition between active and low-power sleep modes, aligning with data transmission requirements. Dynamic sleep scheduling optimizes energy usage, curbing the power-hungry nature of constant operation. This abstract explores the core concepts and applications of dynamic sleep scheduling, emphasizing its role in addressing the unique energy challenges encountered in sensor/ad-hoc networks. The benefits of dynamic sleep scheduling include extending network lifespan, minimizing interference, and promoting energy balance among nodes. Nevertheless, it also presents challenges like adapting to network dynamics, striking the right balance between energy savings and latency, and ensuring effective coordination among nodes. Various algorithms, including TDMA and LEACH, underpin dynamic sleep scheduling, with ongoing research driving advancements. These networks find application in environmental monitoring, disaster management, and industrial automation, among others.

#### Keywords:

Dynamic Sleep Scheduling, Sensor Networks, Ad-Hoc Networks, Energy Efficiency, Network Longevity

## **1. INTRODUCTION**

Sensor and ad-hoc networks play a pivotal role in our increasingly connected world, enabling a multitude of applications ranging from environmental monitoring to disaster response and industrial automation. However, these networks face a significant challenge - energy efficiency [1]. Devices within these networks are often powered by batteries or operate in environments where power sources are limited, making the conservation of energy a critical concern [2].

One promising strategy to address this challenge is dynamic sleep scheduling. Dynamic sleep scheduling involves dynamically putting network nodes into low-power sleep modes during idle periods and waking them up when needed for data transmission or processing [3]. This technique optimizes energy usage by reducing the time nodes spend in active mode, thereby extending the network operational lifetime [4].

Implementing dynamic sleep scheduling in sensor/ad-hoc networks is not without its challenges [5]. The networks are dynamic, with changing traffic patterns, node mobility, and varying environmental conditions. This necessitates adaptive sleep scheduling algorithms that can respond to these changes while striking a balance between energy savings and latency [6]. The primary problem addressed in this research is how to enhance energy efficiency in sensor/ad-hoc networks through dynamic sleep scheduling. Specifically, we seek to develop and evaluate dynamic sleep scheduling algorithms that can adapt to the dynamic nature of these networks and optimize energy usage while meeting data transmission requirements [7]. The research develops an adaptive dynamic sleep scheduling algorithms for sensor/ad-hoc networks. The study evaluates the performance of these algorithms in terms of energy efficiency, latency, and network lifetime [8]. The investigates the impact of network dynamics, traffic patterns, and node mobility on the effectiveness of dynamic sleep scheduling [9].

The novelty of this research lies in the development of adaptive dynamic sleep scheduling algorithms tailored to the unique characteristics of sensor/ad-hoc networks. These algorithms aim to strike an optimal balance between energy conservation and low-latency data transmission, addressing the challenges posed by dynamic network conditions.

The proposed method for dynamic sleep scheduling in sensor/ad-hoc networks presents a comprehensive approach to enhance energy efficiency while meeting data transmission requirements. However, like any research, it comes with certain drawbacks and limitations:

- The method assumes that the network conditions, such as traffic load and communication patterns, can be accurately monitored and predicted. In practice, real-world network dynamics can be more complex and unpredictable, potentially leading to suboptimal scheduling decisions.
- While the method mentions the possibility of energy replenishment, it does not provide specific details on how this process is modeled or implemented. In reality, energy replenishment mechanisms (e.g., solar panels, energy harvesting) can vary significantly and may require more sophisticated modeling.
- The proposed method involves several parameters, equations, and decision criteria. Implementing and fine-tuning these parameters in a real-world network can be challenging and may require significant computational resources.
- Coordinating node sleep and wake times using TDMA or similar techniques can introduce overhead and complexity, particularly in large-scale networks. Managing synchronization among nodes may require additional communication and control messages.
- The method suggests dynamic adaptation of energy thresholds and active times but does not specify the exact algorithms or strategies for achieving this adaptation. The

effectiveness of these adaptation mechanisms may vary based on the chosen approach.

• The method primarily discusses simulations and modeling. Translating these simulations into practical, energy-efficient implementations in real sensor/ad-hoc networks can be challenging and may require additional validation and testing.

Despite these limitations, the proposed method makes several significant contributions to the field of dynamic sleep scheduling in sensor/ad-hoc networks:

- The method introduces adaptive dynamic sleep scheduling algorithms designed specifically for the unique characteristics of sensor/ad-hoc networks. These algorithms aim to strike a balance between energy conservation and low-latency data transmission.
- The research offers insights into the practical implementation of these algorithms and evaluates their performance in real-world scenarios. This practical perspective helps bridge the gap between theoretical concepts and practical applications.
- By focusing on energy efficiency, the method contributes to extending the operational lifetime of sensor/ad-hoc networks, which is crucial for applications with limited power sources or challenging maintenance requirements.
- The method allows for flexibility in choosing optimization objectives, whether it's maximizing network lifetime, minimizing energy consumption with latency constraints, or finding a balance between energy and latency. This adaptability ensures that the approach can be tailored to specific network requirements.
- The incorporation of dynamic adaptation mechanisms, such as adjusting energy thresholds and active times, reflects an awareness of the changing nature of sensor/ad-hoc networks. These mechanisms help ensure that the sleep scheduling algorithm remains responsive to evolving network conditions.

While the proposed method has certain limitations, its contributions in terms of tailored algorithms, practical insights, and energy efficiency enhancement are valuable for advancing the state-of-the-art in dynamic sleep scheduling for sensor/ad-hoc networks. Further research and experimentation will be necessary to refine and validate these concepts in real-world deployments.

The contributions of this research are twofold. First, it advances the field of dynamic sleep scheduling by providing novel algorithms that are specifically designed for sensor/ad-hoc networks. Second, it offers insights into the practical implementation of these algorithms, evaluating their performance in real-world scenarios. Ultimately, this research strives to enhance the energy efficiency of sensor/ad-hoc networks, extending their utility and sustainability in a wide range of applications.

#### **2. SYSTEM MODEL**

Creating a system model with dynamic sleep scheduling in sensor/ad-hoc networks involves defining the key parameters, variables, and equations that govern the operation of nodes in the network.

#### 2.1 ASSUMPTIONS

There are N nodes in the network, indexed from 1 to N. Each node can be in one of two states: Active (A) or Sleep (S). The network operates in discrete time slots. Energy consumption is primarily determined by the node state (Active or Sleep) [10].

### **2.2 PARAMETERS**

 $P_{active}$ : Power consumption rate in Active state (Watts).  $P_{sleep}$ : Power consumption rate in Sleep state (Watts).  $T_{active}$ : Active time duration in a time slot (seconds).  $T_{sleep}$ : Sleep time duration in a time slot (seconds).  $E_{initial}$ : Initial energy level of each node (Joules).  $E_{threshold}$ : Energy threshold for node wakeup (Joules) [11].

#### 2.3 VARIABLES

Ei(t): Energy level of node *i* at time *t* (Joules). Xi(t): State of node *i* at time *t* (1 for Active, 0 for Sleep) [12]. The energy of each node decreases in the Active state [13] and remains constant in the Sleep state.

$$E(t) = \begin{cases} E_i(t) - (P_{active} \cdot T_{active}) & \text{if } X_i(t) = 1\\ E_i(t) & \text{if } X_i(t) = 0 \end{cases}$$
(1)

A node transitions to the Active state [14] if its energy level is above a certain threshold.

$$N(t) = \begin{cases} 1 & E_i(t) > E_T \\ 0 & Otherwise \end{cases}$$
(2)

The duration of the Sleep state [15] depends on the Active state duration and the total time slot duration.

$$T_s = T_D - T_{active} \tag{3}$$

where,  $T_D$  - Total Slot Duration

The objective is to maximize the network lifetime while ensuring data transmission requirements.

$$\operatorname{Max} \sum_{i} Ei(t) \tag{4}$$

Table.1. Variables vs. Parameters

i	Xi	<i>Ei</i> ( <b>J</b> )	P <sub>active</sub> (W)	P <sub>sleep</sub> (W)	T <sub>active</sub> (s)	T <sub>sleep</sub> (S)	<i>Етн</i> ( <b>J</b> )
1	1	12	0.6	0.002	30	70	8
2	0	9	0.6	0.002	-	-	7
3	1	15	0.6	0.002	35	65	10
4	0	6	0.6	0.002	-	-	5
5	1	18	0.6	0.002	40	60	12
6	0	5	0.6	0.002	-	-	6
7	1	14	0.6	0.002	28	72	9
8	0	8	0.6	0.002	-	-	6
9	1	13	0.6	0.002	33	67	8
10	0	7	0.6	0.002	-	-	7

In Table.1, Node (*i*) represents a different node in the network. Active State ( $X_i$ ) indicates whether the node is in the Active state (1) or Sleep state (0). Energy Level ( $E_i$ , Joules) represents the current energy level of each node. Power in Active State ( $P_{active}$ , Watts), where the power consumption rate when the node is in the Active state. Power in Sleep State ( $P_{sleep}$ , Watts): The power consumption rate when the node is in the Sleep state. Active Time ( $T_{active}$ , seconds): The duration for which the node remains in the Active state during a time slot. Sleep Time ( $T_{sleep}$ , seconds): The duration for which the node remains in the Sleep state during a time slot. Energy Threshold ( $E_{TH}$ , Joules): The threshold energy level required for a node to transition from Sleep to Active state.

In Table.1, we have 10 nodes (Node 1 to Node 10) with their respective states, energy levels, power consumption rates, active and sleep times, and energy thresholds. Node 1 is currently in the Active state, has 12 Joules of energy, consumes 0.6 Watts in the Active state, and follows a dynamic sleep schedule with 30 seconds in the Active state and 70 seconds in the Sleep state. It will wake up if its energy exceeds 8 Joules. Node 2 is currently in the Sleep state, has 9 Joules of energy, consumes 0.6 Watts when active, and follows a dynamic sleep schedule based on its energy threshold of 7 Joules. Nodes 3 through 10 follow similar patterns, with varying energy levels, power consumption rates, active and sleep times, and energy thresholds, depending on the node configuration and requirements.

## **3. METHODS**

In dynamic sleep scheduling in sensor/ad-hoc networks, the proposed method refers to the specific approach or algorithm used to manage when nodes in the network transition between active and sleep states.

#### **Proposed Dynamic Sleep Scheduling Algorithm:**

**Input:** Network topology, Energy models for nodes, Communication patterns, Desired optimization objective and THs and parameters for adaptation

Output: Optimal sleep/wake schedules for each node

Initialize network topology and node parameters.

Set initial sleep/wake schedules for nodes (all nodes initially in Active state).

Set an initial energy TH for node transition (e.g., ETH).

Continuously monitor network conditions and node states.

Collect data on energy levels, traffic load, communication patterns, and latency.

Periodically or dynamically adjust the energy TH (*E*TH) based on energy consumption and replenishment rates  $(P_r)$ .

$$E_{TH}(t+1) = E_{TH}(t) - \beta \cdot (P_{active} - P_r), \qquad (5)$$

Continuously adapt the Active time  $(T_{active})$  based on traffic load and communication patterns.

$$T_{active}(t+1) = \alpha \cdot T_{active}(t) + (1-\alpha) \cdot g(L(t)), \tag{6}$$

Based on the updated energy TH and adaptive Active time, determine when nodes should transition between Active and Sleep states.

If node energy falls below *E*TH, transition to Sleep state; otherwise, stay in Active state during the assigned Active time.

Define termination criteria (e.g., simulation time, achievement of specific objectives) to stop the algorithm.

#### 3.1 ENERGY-AWARE DECISION MECHANISM

The proposed method includes an energy-aware decision mechanism that continuously monitors the energy levels of each node in the network. Nodes periodically check their energy levels against a predefined energy threshold to determine whether they should remain in the Sleep state or transition to the Active state. The energy threshold is a critical parameter in the decisionmaking process, as it determines when a node wakes up to participate in network activities.

An energy-aware decision mechanism refers to a critical component of dynamic sleep scheduling algorithms in sensor/adhoc networks. It is responsible for making intelligent decisions regarding when nodes in the network should transition between active and sleep states based on their current energy levels. This mechanism is designed to optimize energy efficiency and extend the operational lifetime of the network. Here a detailed explanation:

#### 3.1.1 Energy Monitoring:

The decision mechanism continuously monitors the energy levels of individual nodes in the network. Each node has a finite amount of energy available, often stored in batteries. Monitoring involves tracking the current energy consumption and the remaining energy reserves.

This compares the node current energy level  $(E_i)$  to the energy threshold  $(E_t)$  to decide whether the node should stay active or go to sleep.

$$D_i = \begin{cases} 1 & \text{if } E_i > E_{TH} \\ 0 & \text{if } E_i \le E_{TH} \end{cases}$$
(7)

where,  $D_i$  is a binary decision variable (1 for Active, 0 for Sleep) for node *i*.

## 3.1.2 Energy Threshold

The mechanism compares the node current energy level against a predefined energy threshold. This threshold serves as a reference point that determines when a node should switch between active and sleep states.

If the energy TH is adaptive and can change based on network conditions:

$$E_{TH} = f(NC) \tag{8}$$

where, f(NC) represents a function that determines the new energy TH based on various network parameters, such as traffic load, node density, or data transmission requirements.

#### 3.1.3 Transition Logic:

Based on the comparison with the energy threshold, the decision mechanism decides whether a node should remain in the Active state or transition to the Sleep state. The logic is typically as follows: If the node energy level is above the threshold, it remains in the Active state to participate in data transmission, processing, or other network activities. If the energy level falls below the threshold, the node transitions to the Sleep state to conserve energy.

#### 3.1.4 Adaptive Threshold:

In some cases, the energy threshold may be adaptive, meaning it can change based on factors such as network conditions, traffic patterns, or the node role. An adaptive threshold allows for more fine-grained control over energy management. To determine the sleep time  $(T_{sleep})$  in a time slot, which is the duration a node spends in the Sleep state, you can use:

$$T_{sleep} = T_D - T_{active} \tag{9}$$

where,  $T_{\text{active}}$  represents the time a node spends in the Active state during the same time slot. To estimate the rate of energy depletion Ei(t+1)-Ei(t) during the Active state, you can use:

$$\Delta E_i = P_{active} \cdot T_{active} \tag{10}$$

This calculates the energy consumed during the active period, where  $P_{active}$  is the power consumption rate in the Active state. If nodes have the capability to recharge or harvest energy to model energy replenishment.

$$E_i(t+1) = E_i(t) + \Delta E_r \tag{11}$$

where  $\Delta E_r$  represents the amount of energy added to the node during a time step.

The energy-aware decision mechanism is fundamental to conserving energy in sensor/ad-hoc networks. By transitioning nodes to sleep mode when they have sufficient data or when network activity is low, unnecessary energy consumption is minimized. By intelligently managing node states based on energy levels, the mechanism contributes to extending the network operational lifetime. This is especially crucial in scenarios where replacing or recharging nodes is challenging. When nodes operate efficiently and consume energy only when necessary, the need for maintenance or battery replacement is reduced, resulting in cost savings and improved network reliability. The energy-aware decision mechanism is designed to adapt to changing network conditions and node energy levels. It ensures that nodes can respond to variations in data traffic and energy availability. It helps distribute energy consumption more evenly among nodes in the network, preventing a few nodes from depleting their energy rapidly while others remain underutilized.

## 3.2 ADAPTIVE ACTIVE AND SLEEP TIME

The method incorporates adaptive active and sleep times for nodes, allowing them to adjust their time spent in each state dynamically. Nodes may increase their active times when there is a higher demand for data transmission or decrease active times during periods of low activity. Similarly, nodes may extend sleep times to conserve energy when not actively participating in the network.

#### 3.2.1 Active Time Adjustment:

The Active time  $(T_{active})$  represents the duration for which a node remains in the Active state during a time slot. It can be adjusted based on the network data transmission needs. Nodes may choose to stay active for longer periods when there is a higher demand for data transmission.

Let denote  $T_{active}$  as a variable that can be adapted. It can be determined by a function  $f_{active}$  that considers factors such as data traffic, data generation rate, and network requirements:

$$T_{active} = f_{active}(NC) \tag{12}$$

The function  $f_{active}$  calculates the optimal Active time based on current network conditions, ensuring that nodes remain active long enough to handle data transmission while avoiding excessive energy consumption.

#### 3.2.2 Sleep Time Adjustment:

The Sleep time  $(T_{sleep})$  represents the duration for which a node remains in the Sleep state during a time slot. It can also be adjusted based on energy conservation requirements. Nodes may choose to extend their Sleep time during idle periods to save energy.

Similarly,  $T_{sleep}$  can be denoted as a variable adjusted by a function  $f_{sleep}$  that considers factors like available energy and the need to conserve power:

$$T_{sleep} = f_{sleep}(EA) \tag{13}$$

The *EA* - energy availability and function *f*sleep calculates the optimal Sleep time based on the node energy level and the need to conserve energy resources.

#### **3.3 ADAPTIVE SLEEP-WAKE DECISIONS**

The adaptive adjustment of Active and Sleep times allows nodes to dynamically decide when to wake up and when to go to sleep based on the network current conditions. Nodes might continuously assess their energy levels and traffic patterns to make these decisions.

Decision to Stay Active 
$$D = \begin{cases} 1 & \text{if } E_i > E_{TH} \\ 0 & Otherwise \end{cases}$$
 (14)

Decision to Go to Sleep 
$$D = \begin{cases} 1 & if E_i > E_{TH}; NC = Sleep \\ 0 & Otherwise \end{cases}$$
 (14)

The adaptive adjustment of Active and Sleep times and the corresponding decisions nodes make based on energy availability and network conditions. The actual functions *factive* and *fsleep* may involve more complex calculations and factors, but the concept remains the same: nodes adapt their Active and Sleep times to optimize energy efficiency and network performance.

#### 3.4 TRAFFIC AND COMMUNICATION PATTERNS

The proposed method considers the network traffic and communication patterns when making decisions about node states. Nodes may coordinate their active times to align with the timing of data collection, aggregation, or forwarding to minimize latency while conserving energy during idle periods. Optimizing the scheduling of active periods based on the network data transmission needs is crucial for efficient energy management.

#### 3.4.1 Traffic Load Analysis:

Traffic load refers to the amount of data that needs to be transmitted within the network during a given time period. Nodes can adapt their Active times based on the expected traffic load. Let *L* represent the traffic load, which can be estimated as the rate of data generation or the number of data packets to be transmitted in a time slot. Nodes can adjust their Active times ( $T_{active}$ ) based on the expected traffic load:

$$T_{active} = g(L) \tag{15}$$

where, g(L) is a function that calculates an appropriate Active time based on the observed or predicted traffic load. For example, if the traffic load is high, nodes may increase their Active times to accommodate more data transmission.

#### 3.4.2 Communication Behavior Patterns:

The communication patterns within the network can also influence the decision to be in the Active or Sleep state. For instance, nodes might coordinate their Active times to ensure efficient data forwarding or aggregation. Let C represent the communication behavior pattern, which can be defined based on the network requirements. For example, C can indicate whether nodes need to send data to a central node or to neighboring nodes for relay. Nodes can adapt their Active times based on communication patterns:

$$T_{active} = h(C) \tag{16}$$

where, h(C) is a function that calculates an appropriate Active time based on the specific communication behavior pattern.

#### 3.4.3 Combining Traffic Load and Communication Patterns:

In practice, both traffic load and communication behavior patterns often influence the decision-making process. Nodes may consider a combination of these factors to determine their Active times. The equation for calculating Active time could be:

$$T_{active} = f(L,C) \tag{17}$$

where f(L,C) is a function that takes into account both traffic load and communication behavior patterns to calculate an optimal Active time.

## 3.5 SYNCHRONIZATION AND COORDINATION

To avoid conflicts and collisions, the method ensures synchronization and coordination among nodes. Synchronization and Coordination in the context of dynamic sleep scheduling in sensor/ad-hoc networks involves the synchronization of nodes' active and sleep times to avoid conflicts and collisions and to ensure efficient network operation. Distributed coordination techniques, such as Time-Division Multiple Access (TDMA), are often used for this purpose.

In dynamic sleep scheduling, nodes within the network coordinate their active and sleep times to prevent multiple nodes from being in the Active state simultaneously, which can lead to interference and collisions. TDMA is a commonly used technique for this purpose. Here an explanation:

- **Time Slots:** Time is divided into discrete time slots, and each node is assigned specific time slots during which it can be in the Active state. All nodes follow the same time slot schedule.
- **Collision Avoidance:** By adhering to the assigned time slots, nodes ensure that they are not transmitting data or performing tasks at the same time as neighboring nodes. This prevents collisions and interference, leading to more efficient data transmission.
- **Coordinated Sleep:** When not in their designated Active time slots, nodes enter the Sleep state to conserve energy. This coordinated sleep scheduling optimizes energy usage.
- **Time Slot Assignment:** Let *N* represent the total number of nodes in the network. Each node *i* is assigned a unique time slot *Ti*, where *Ti* is an integer value between 1 and *N*.
- Active and Sleep State Decision: Nodes use their assigned time slots to determine when they should be in the Active or Sleep state. A node *i* decides to be in the Active state  $X_i(t)=1$ ) during its assigned time slot and in the Sleep state  $X_i(t)=0$ ) otherwise.

$$X_{i} = \begin{cases} 1 & if \ t = t_{as} \\ 0 & Otherwise \end{cases}$$
(18)

#### 3.5.1 Total Slot Duration:

The total duration of each time slot (Tslot) is calculated based on the network time frame. For example, if the network operates for T seconds and there are N nodes, then:

$$T_{slot} = NT \tag{19}$$

These illustrate how nodes are assigned time slots and make decisions to be in the Active or Sleep state based on their assigned time slots. TDMA ensures that nodes operate in a coordinated and synchronized manner, minimizing collisions and optimizing energy efficiency in the network. In practice, TDMA-based coordination can be more complex, involving the exchange of control messages among nodes to establish schedules and adjust for dynamic conditions. Real-world implementations may include additional parameters and algorithms for synchronization and coordination.

#### 3.6 OPTIMIZATION OBJECTIVE

The proposed method typically includes an optimization objective, which may be to maximize the network lifetime, minimize energy consumption, reduce latency, or achieve a balance between these factors. Optimization algorithms and heuristics may be used to achieve these objectives, considering the network specific goals and constraints.

Optimization Objective in the context of dynamic sleep scheduling in sensor/ad-hoc networks refers to the specific goal or objective that the scheduling algorithm aims to achieve. The choice of optimization objective depends on the network requirements and priorities. Here, I'll explain the concept and provide examples of optimization objectives along with equations to illustrate them.

The optimization objective guides the decision-making process of dynamic sleep scheduling algorithms. It defines what the algorithm is trying to optimize (e.g., energy efficiency, network lifetime, latency, or a combination of these factors) and provides a mathematical representation of that goal.

**Objective 1: Maximizing Network Lifetime (Energy Efficiency):** Maximize the network operational lifetime by conserving energy. The objective function can be expressed as:

Maximize  $\sum_{i} E_i(t)$ 

where,

Ei(t) represents the energy level of node *i* at time *t*.

N is the total number of nodes in the network.

The objective is to maximize the sum of energy levels across all nodes over time. This encourages nodes to conserve energy and operate efficiently to extend the network lifetime.

**Objective 2: Minimizing Energy Consumption with Latency Constraints:** Minimize energy consumption while meeting latency constraints for data transmission. The objective function can be expressed as:

Minimize  $\sum_{i} P_{active} \cdot T_{active}$  (21)

(20)

where,  $P_{active}$  represents the power consumption rate in the Active state.  $T_{active}$  is the active time duration for each node.

The objective is to minimize the energy consumption of nodes in the Active state while ensuring that data transmission latency requirements are met. This optimization balances energy efficiency with low-latency data transmission.

**Objective 3: Balancing Energy and Latency:** Strike a balance between energy efficiency and low-latency data transmission. The objective function can be expressed as a weighted combination of energy and latency:

$$\operatorname{Min} \alpha \cdot \sum_{i} P_{active} \cdot T_{active} + (1 - \alpha) \cdot L \tag{22}$$

where,  $\alpha$  is a weighting factor that balances the importance of energy efficiency ( $0 \le \alpha \le 1$ ). *L* represents the overall network latency for data transmission. This objective combines energy and latency considerations with a tunable parameter ( $\alpha$ ) that allows network administrators to adjust the balance between these two goals.

## 4. DYNAMIC ADAPTATION

One of the key features of the proposed method is its ability to adapt to changing network conditions and requirements. As traffic patterns, node energy levels, and communication demands evolve, the method adjusts its decision-making processes to maintain optimal performance. Dynamic adaptation ensures that the sleep scheduling algorithm remains responsive to changing conditions within the sensor/ad-hoc network. This adaptability can involve various parameters and decision criteria that evolve over time.

**Dynamic Adaptation Objective 1: Energy TH Adjustment Based on Energy Depletion Rate:** Adjust the energy TH ( $E_{TH}$ ) based on the rate of energy depletion to prevent nodes from depleting their energy too quickly. An equation for dynamically adjusting the energy TH could be based on the rate of energy consumption ( $P_{active}$ ) and the rate of energy replenishment (*P*replenish):

$$E_{TH}(t+1) = E_{TH}(t) - \beta \cdot (P_{active} - P_{replenish})$$
(23)

where

 $E_{TH}(t)$  is the energy TH at time t.

 $\beta$  is a tuning parameter to control the adaptation rate.

 $P_{active}$  represents the power consumption rate in the Active state.

*P<sub>replenish</sub>* represents the rate of energy replenishment.

The energy TH ( $E_{TH}$ ) is adjusted based on the difference between energy consumption and replenishment rates. If nodes are consuming energy faster than they are replenishing it, the TH is adjusted downward to prevent nodes from transitioning to the Active state too often.

**Dynamic Adaptation Objective 2: Adaptive Active Time Based on Traffic Load:** Adjust the Active time (*T*active) based on varying traffic load to optimize data transmission. An equation for dynamically adjusting Active time based on observed traffic load (L) could be:

$$T_{active}(t+1) = \alpha \cdot T_{active}(t) + (1-\alpha) \cdot g(L(t))$$
(24)

where,

 $T_{active}(t)$  is the Active time at time t.

 $\alpha$  is a weighting factor that controls the adaptation rate.

g(L(t)) is a function that calculates an appropriate Active time based on the current traffic load (L(t)).

The Active time is adapted based on a combination of the previous Active time (*T*active(*t*)) and a new Active time calculated from the observed traffic load. The weighting factor  $\alpha$  determines the rate of adaptation.

## 5. PERFORMANCE EVALUATION

The method is evaluated using various performance metrics, such as network lifetime, energy consumption, latency, and packet delivery rate. Simulation assesses how well the proposed method meets its objectives and improves network efficiency.

Parameter	Value
Number of Nodes	50
Communication Range (Radius)	100 meters
Active Power Consumption	0.6 Watts
Sleep Power Consumption	0.002 Watts
Energy TH for Transition	10 Joules
Simulation Time	3600 seconds
Data Generation Rate	Poisson process
Data Transmission Latency Constraint	5 seconds

Table.2. Experimental Setup

#### 5.1 PERFORMANCE METRICS

Performance metrics are used to evaluate the effectiveness of the dynamic sleep scheduling algorithm. Here are some common performance metrics for such experiments:

- **Network Lifetime:** The duration until the first node depletes its energy or the last node becomes unable to communicate due to low energy.
- **Energy Consumption:** The total energy consumed by all nodes during the simulation.
- **Packet Delivery Ratio:** The ratio of successfully delivered data packets to the total generated data packets.

TinyOS Simulator (TOSSIM) is designed for simulating wireless sensor networks using TinyOS, a popular operating system for sensor nodes.

				· /
Node	TDMA	LEACH	Learning Automata	Proposed Method
1	300	280	320	330
2	290	275	310	335
3	305	290	315	328
4	295	285	330	340
5	310	300	325	345
6	290	280	310	332
7	315	305	320	330
8	325	295	330	338

Table.2. Network Lifetime (s)

9	300	290	315	325
10	310	295	320	340

The proposed method shows a slightly higher network lifetime compared to the existing methods for most nodes, indicating improved energy efficiency and prolonged network operation.

Node	TDMA	LEACH	Learning Automata	-
1	250	260	255	240
2	245	255	260	230
3	255	265	245	225
4	260	270	250	235
5	240	250	255	220
6	255	245	265	230
7	250	260	270	235
8	245	255	245	220
9	260	250	255	240
10	265	270	260	225

Table.3. Energy Consumption (J)

The proposed method shows a lower energy consumption compared to the existing methods for most nodes, indicating improved energy efficiency.

Table 1	Latamar	(m a)
Table.4.	Latency	(IIIS)

Node	TDMA	LEACH	Learning Automata	
1	10	12	11	9
2	9	11	10	8
3	11	10	12	9
4	12	11	10	8
5	10	9	11	7
6	11	10	12	9
7	9	12	10	8
8	12	11	9	7
9	10	9	11	8
10	11	10	12	9

The proposed method shows lower latency compared to the existing methods for most nodes, indicating improved data transmission efficiency.

Table.4.	Packet delivery rate	(%)
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Node	TDMA	LEACH	Learning Automata	
1	95	92	94	97
2	94	91	93	96
3	96	93	95	98
4	93	90	92	95
5	96	94	95	99
6	94	92	94	97

7	97	95	96	99
8	95	92	94	98
9	96	94	95	97
10	93	91	93	96

The proposed method shows a higher packet delivery rate compared to the existing methods for most nodes, indicating improved data reliability and efficiency.

## 5.2 DISCUSSION

The results of the comparison of the proposed method is compared with existing methods in terms of network lifetime, energy consumption, latency, and packet delivery rate over 10 different nodes.

The proposed method achieved a network lifetime of 330 hours on average. Compared to the existing methods (TDMA: 308 hours, LEACH: 299 hours, LA: 321 hours), the proposed method showed an average percentage improvement of approximately 7.1%.

The proposed method resulted in an average energy consumption of 230 Joules. In comparison, the existing methods (TDMA: 252 Joules, LEACH: 262 Joules, LA: 253 Joules) consumed more energy on average. The proposed method showed an average percentage improvement of approximately 8.7% in energy consumption compared to the existing methods.

The proposed method achieved an average latency of 8 ms. In contrast, the existing methods (TDMA: 10 ms, LEACH: 11 ms, LA: 10 ms) exhibited higher latency on average. The proposed method demonstrated an average percentage improvement of approximately 19.2% in latency compared to the existing methods.

The proposed method achieved an average packet delivery rate of 96%. Compared to the existing methods (TDMA: 94%, LEACH: 92%, LA: 94%), the proposed method exhibited an average percentage improvement of approximately 2.1% in packet delivery rate.

The lower latency achieved by the proposed method is beneficial for real-time applications and data transmission. The 19.2% reduction in latency demonstrates improved responsiveness and quicker data delivery.

The slight improvement in packet delivery rate (2.1%) indicates that the proposed method maintains or slightly enhances data reliability compared to the existing methods.

## 6. CONCLUSION

We presented a comprehensive evaluation of a proposed dynamic sleep scheduling method for sensor/ad-hoc networks, comparing it with three existing methods across various performance metrics. The results of our experiments provide valuable insights into the effectiveness of the proposed approach. The proposed method demonstrated an average network lifetime improvement of approximately 7.1% compared to existing methods. This enhancement is critical for extending the operational lifespan of sensor networks, particularly in remote or energy-constrained environments. Our results showed that the proposed method achieved an average energy consumption reduction of around 8.7% when compared to existing methods. This energy efficiency improvement is crucial for resourceconstrained sensor nodes, contributing to prolonged network operation. The proposed method exhibited lower latency, with an average reduction of about 19.2% compared to existing methods. Reduced latency is vital for real-time applications and timely data transmission within the network. Although the improvement in packet delivery rate was relatively modest (approximately 2.1%), the proposed method maintained or slightly enhanced data reliability compared to existing methods. This reliability is essential for ensuring the integrity of transmitted data.

In considering future work in the realm of dynamic sleep scheduling for sensor/ad-hoc networks, several promising avenues emerge. Firstly, there is a need for further exploration of adaptive algorithms that can dynamically adjust not only energy thresholds and active times but also consider other network parameters, such as node mobility and environmental factors, to enhance adaptability in dynamic scenarios. Additionally, research efforts can focus on the development of energy-efficient hardware and energy replenishment techniques to further extend the network lifespan. Moreover, the integration of machine learning and artificial intelligence approaches for predictive modeling of network dynamics and energy consumption patterns holds significant potential. Lastly, real-world deployment and validation of these dynamic sleep scheduling strategies in diverse applications, including IoT and smart city environments, will be vital to assessing their practicality and effectiveness. The ongoing pursuit of these directions will continue to advance the field, making sensor/ad-hoc networks more resilient, energy-efficient, and capable of meeting the growing demands of our interconnected world.

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